Introduction to the Special Issue on AI & Networks

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As networks have permeated our world, the economy has come to resemble an ecology of organisms, interlinked and coevolving, constantly in flux, deeply tangled, ever expanding at its edges.

Kevin Kelley

Most people . . . would agree that a fundamental property of complex systems is that they are composed of a large number of components or "agents," interacting in some way such that their collective behavior is not a simple combination of their individual behaviors.

Mark Newman

The importance of networks permeates the world today. From biology to social systems, from the brain to the Internet, networks play an important and central role in the way the world works. In the last ten years, due in part to large increases in computational power, large-scale, real-world networks have received much attention from a variety of fields of study.

Within the artificial intelligence community, networks appear some form in nearly every subdiscipline: knowledge representation, inference, learning, natural language processing, multiagent systems, analogical reasoning, and many others. The goals of this special issue are to provide a sampling of research efforts focused on how networks can be used in AI systems, and to facilitate cross-communication among subdisciplines that are studying networks from different perspectives.

The seven papers we include here cover a broad range of network-inspired AI research—in natural language processing, data mining, the Semantic Web, peer-to-peer networks, multi-agent systems, analog networks, and the modern social network of the "blogophere." Each article represents a snapshot of the area it describes; for example, the collective classification problem surveyed by Sen et al. is just one of many problems within the emerging research area of link mining. Moreover, networks are influential in many other areas of AI that are not represented here, including Bayesian networks and graphical models, sensor networks, swarm systems and cellular automata, graphical games, trust and reputation systems, and computational organizational design, just to name a few.

Table 1 summarizes the articles in this collection by characterizing the nature of the networks that are the focus of each of the seven papers.

1 Basic Graph Theory

In reading the articles presented here, some basics of graph and network theory may be useful for the reader who is not familiar with these terms. We start with some basic terminology:

- A graph G is defined to be a pair (V, E), where V is a vertex set and E is an edge set (see below). The terms graph and **network** are often used interchangeably.
- The finite vertex set V is a set of descriptors that describe the vertices in the graph. Each vertex may just have an identifier, or it may have an arbitrarily complex set of attributes. The terms vertex and node are often used interchangeably. Depending on the application, nodes may also be referred to as agents or entities.

Authors	Topic	Nodes	Edges	Tasks
Radev &	Natural	Words, word	Co-occurrences;	Analyze syntax; iden-
Mihalcea	language	senses,	collocations;	tify lexical semantics; re-
	processing	sentences,	syntactic structure;	trieve and summarize
		documents	lexical similarity	text; extract keywords
Berners-Lee &	Semantic	Agents,	Connections	Disseminate knowledge;
Kagal	Web	terms,	between	construct and share on-
		ontologies	communities;	tologies; provide and re-
			subtask	quest services; create
			relationships;	new communities
			ontological	
			relationships	
Menczer,	Peer-to-peer	Agents	"Social" connections	Locate relevant knowl-
Wu, &	networks		along which queries	edge sources; learn which
Akavipat			flow	peers can answer queries
Pearce,	Cooperative	Agents	Interactions, joint	Multi-agent plan co-
Tambe, &	multi-agent		reward structures	ordination, meeting
Maheswaran	systems			scheduling, teamwork
				(e.g., RoboCup soccer)
Mattiussi,	Analog	Dynamical	Signal flows with	Synthesize and reverse-
Mrabach,	networks	devices	varying strength	engineer analog networks
Dürr, &				(e.g., gene regulatory
Floreano				networks and analog
				electronic circuits)
Finin,	Blogosphere	Web pages,	Social networks;	Recognize spam blogs
Joshi,		blog postings,	comments;	(splogs); find opinions on
Kolari,		bloggers, blog	trackbacks;	topics; identify commu-
Java,		sites	advertisements;	nities of interest; derive
Kale, &			tags; RDF data;	trust relationships; de-
Karandikar			metadata	tect influential bloggers
Sen,	Social and	Entities (e.g.,	Relationships	Perform collective clas-
Namata,	natural	scientific	among the entities	sification; construct fea-
Bilgic,	networks	articles)	(e.g., citations or	tures for relational clas-
Getoor,			co-citations)	sification
Gallagher, &				
Eliassi-Rad				

Table 1: Network types

- The finite edge set E specifies the relationships between the vertices in the graph. Each edge $e \in E$ is a pair of vertices, which are called the **endpoints** of the edge. Edges may be **ordered** or **unordered** and also **weighted** or **unweighted**. A **hyperedge** may connect more than two vertices. Edges are often used to represent **relations**.
- The degree of a node, k_i , is the number of edges that are connected to node i. In directed graphs, degree can be broken down into "in-degree" (number of edges coming into the node) and "out-degree" (number of edges pointing out of the node).

2 Network properties

A number of properties prove to be useful in graph theory and social network theory for analyzing and understanding the behavior of graph structures.

• The **path length** between two nodes is the minimum number of edges that must be traversed to move from one node to the other in the graph. The **average path length** is an average across all pairs of nodes in the graph.

Real-world graphs often exhibit **short average path lengths**, meaning that the average path length is less than would be expected in a random graph. This "small-world effect" was first recognized by Milgram [5] in analyzing the number of hops it took for human subjects to send a piece of postal mail to a predefined destination by following only links to people whom they knew on a first name basis. This phenomenon is sometimes called "six degrees of separation," based on the hypothesis that any two people in the world can be connected by a six-link "chain" of acquaintances. A game created in the mid-nineties called "Six degrees of Kevin Bacon" (find a short path connecting any given movie actor or actress to Kevin Bacon) in fact initiated some of the research work that led to the current boom in interest in network studies.

Several other properties are related to path length:

- The **betweenness** of a node i is the number of other pairs of nodes (j, k) whose shortest paths pass through i.
- The **closeness** of a node is the average shortest path to all other nodes in the graph.
- The **diameter** of a graph is the length of the longest of all shortest paths (i.e., it is the maximal distance between any pair of nodes (i, j)).
- Clustering measures are used to characterize the frequency of transitive relationships in networks [6, 1, 8]. The clustering coefficient of a network is the ratio of triangles in a network (sets of three nodes that are all connected to each other) to the number of connected triples (sets of three nodes in which at least one node is connected to the other two).

Real-world networks often exhibit **excess clustering**, in the sense that they have a much higher (often 2 orders of magnitude or more) clustering coefficient than would be expected in a random graph of the same size [6]. This is because in many processes that generate networks, two nodes that are connected to a common neighbor are more likely to become connected.

• The degree of a node is also sometimes called its degree centrality, since the number of edges that are connected to a node give an indication of how "central" it is to the network. The degree distribution of a network is the frequency of occurrence of nodes with each degree. A useful summary property is the network's average degree, which can be thought of as the density of the network. The normalized standard deviation of the degrees of the nodes can be used to characterize how much variability there is in the network density. The degree correlation of adjacent nodes in a network indicates whether neighboring nodes are likely to have similar degree.

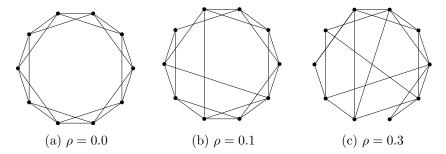


Figure 1: Three increasingly random small-world networks: (a) a small world with no shortcut links; (b) the same small world with a few shortcuts; and (c) a small world with many shortcuts, which begins to resemble a random graph. All three of the networks are constructed from a one-dimensional lattice where nodes are connected to K=2 other nodes in each direction, based on physical proximity. This particular choice of initial layout is deliberate in that it ensures a high initial clustering coefficient.

Many real-world networks have highly skewed degree distributions, with high normalized standard deviation. In particular, the degree distribution in real-world networks often follow a **power law**, where the probability of a node in the network having degree k is proportional to $k^{-\gamma}$ for some parameter γ [6] (typically \$gamma\$ is between -2 and -3). Such networks have a hub-and-spoke structure, with some nodes having very large degree [1].

3 Network models

A variety of network models have been proposed to represent various types of network formation processes and graph behaviors. Several of the most common models are described below.

• Regular graphs have a homogeneous connectivity pattern for all of the nodes in the graph. In these graphs, the degree distribution is trivial: all nodes have the same degree. Examples of regular graphs include lattices, hyper-cubes, and fully connected networks (in which all nodes are connected to all other nodes).

The **coordination number** [7] of a lattice graph determines the number of connections that each node has with its spatial "nearest neighbors" in each dimension. An example of a one-dimensional lattice with K=2 is shown in Figure 1(a).

• Random graphs were first introduced by Erdős and Rényi [4]. A random graph $G_{n,p}$ consists of n nodes where p denotes the probability of an edge existing between each pair of vertices. Random graph models have been widely studied, in part because their properties can be computed analytically. For instance, the expected number of undirected edges in $G_{n,p}$ is n(n-1)p/2, and the average degree of a vertex is k = p(n-1).

A random geometric graph is a special case of a random graph that is generated by randomly placing N agents in the unit square, then connecting pairs of agents if they are within some specified distance d of each other [3]. More specifically, two agents, i and j are connected in a random geometric graph if $d(i,j) < \phi$, where ϕ is a threshold parameter of the model. Figure 2 shows an instance of a random geometric graph with $\phi = 0.09$.

• The **small-world network** model of Watts and Strogatz [8] is an attempt to produce networks that exhibit the real-world properties of excess clustering and short average path length. Small-world networks have properties that lie between those of regular (lattice) networks and random graphs.

Small-world networks are constructed by randomly "re-wiring" each edge in a lattice network with some probability ρ . This process results in shortcut connections across the

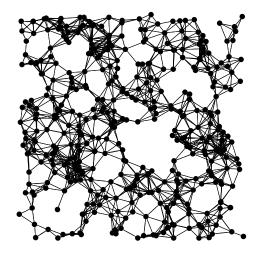


Figure 2: An instance of a random geometric graph on the unit square with 400 nodes and $\phi = 0.09$.

network, as seen in Figure 1. (When edges are replaced with random shortcuts with probability $\rho = 1$, the resulting graph is a random graph.)

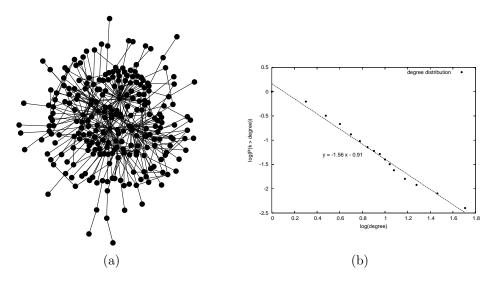


Figure 3: An example of a scale-free network structure with 250 nodes: (a) a rendering of the network that clearly shows the hub-and-spoke structure, and (b) a log-log plot of the cumulative degree distribution of the network shown in (a). Note that a linear curve in a log-log plot implies a power-law behavior of the underlying system.

• The scale-free graph model is motivated by the empirically measured degree distributions of the Internet and the World Wide Web (WWW) [1, 2]. The model is a highly intuitive model based on the way that many networks are believed to evolve and grow in the real world.

The generation of scale-free graphs has two simple rules:

- 1. growth: at each time step, a new node is added to the graph, and
- 2. **preferential attachment:** when a new node is added to the graph, it attaches preferentially to existing nodes with high degree.

Figure 3 shows an example of a scale-free network structure, and the power-law degree distribution that it exhibits.

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